



Instruction Selection

COMP 412
Fall 2005

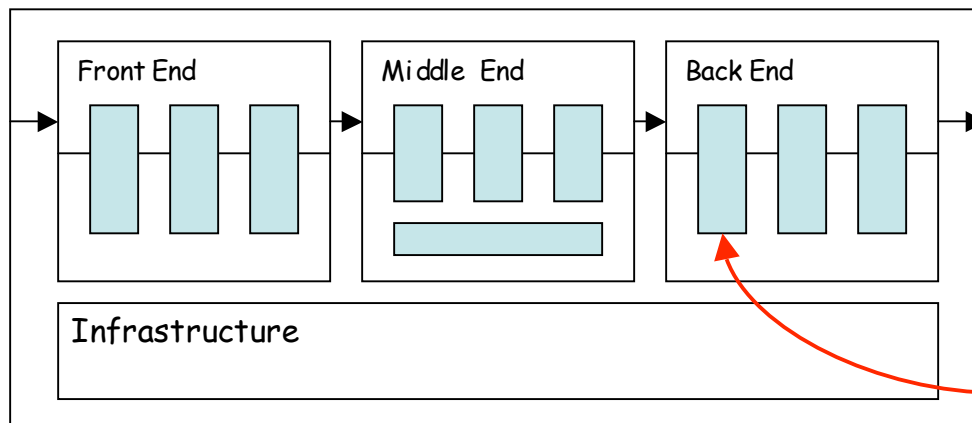
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The Problem

Writing a compiler is a lot of work

- Would like to reuse components whenever possible
- Would like to automate construction of components



Today's lecture:
Automating
Instruction
Selection

- Front end construction is largely automated
- Middle is largely hand crafted
- (Parts of) back end can be automated

Definitions



Instruction selection

- Mapping IR into assembly code
- Assumes a fixed storage mapping & code shape
- Combining operations, using address modes

Instruction scheduling

- Reordering operations to hide latencies
- Assumes a fixed program (*set of operations*)
- Changes demand for registers

Register allocation

- Deciding which values will reside in registers
- Changes the storage mapping, may add false sharing
- Concerns about placement of data & memory operations



The Problem

Modern computers (still) have many ways to do anything

Consider register-to-register copy in ILOC

- Obvious operation is `i2i ri ⇒ rj`
- Many others exist

<code>addI r_i,0 ⇒ r_j</code>	<code>subI r_i,0 ⇒ r_j</code>	<code>lshiftI r_i,0 ⇒ r_j</code>
<code>multI r_i,1 ⇒ r_j</code>	<code>divI r_i,1 ⇒ r_j</code>	<code>rshiftI r_i,0 ⇒ r_j</code>
<code>orI r_i,0 ⇒ r_j</code>	<code>xorI r_i,0 ⇒ r_j</code>	... and others ...

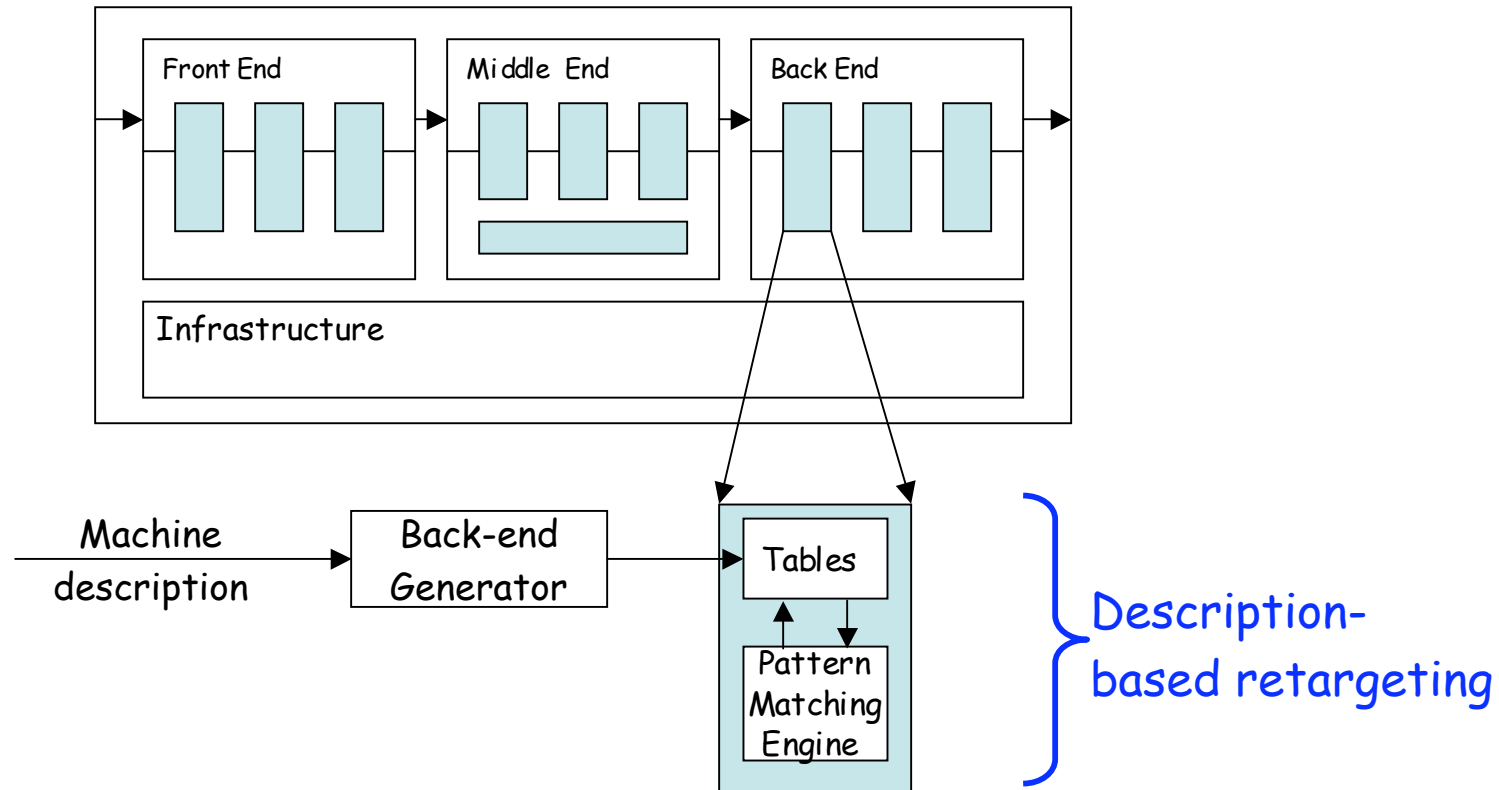
- Human would ignore all of these
- Algorithm must look at all of them & find low-cost encoding
 - Take context into account *(busy functional unit?)*

And ILOC is an overly-simplified case



The Goal

Want to automate generation of instruction selectors



Machine description should also help with scheduling & allocation



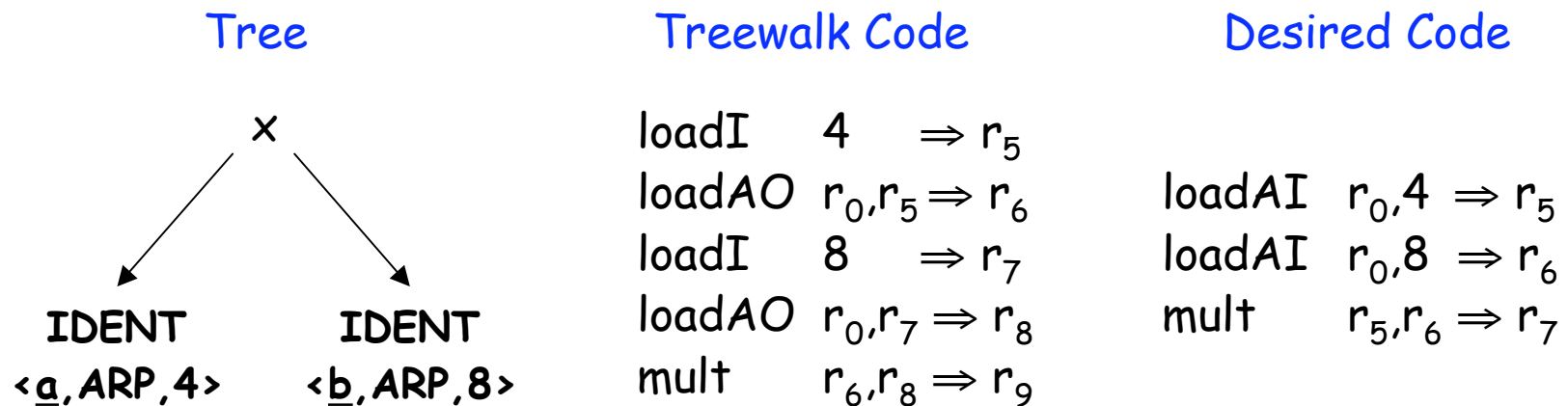
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator (Lec. 24) ran quickly

How good was the code?





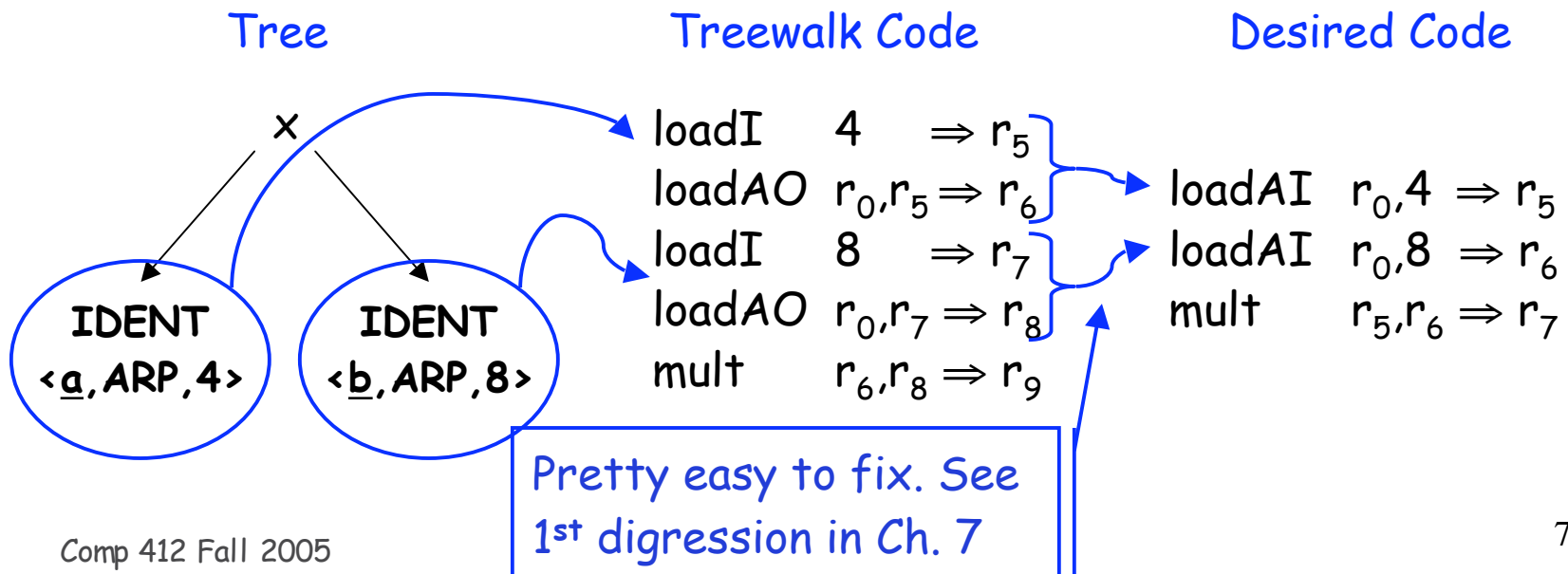
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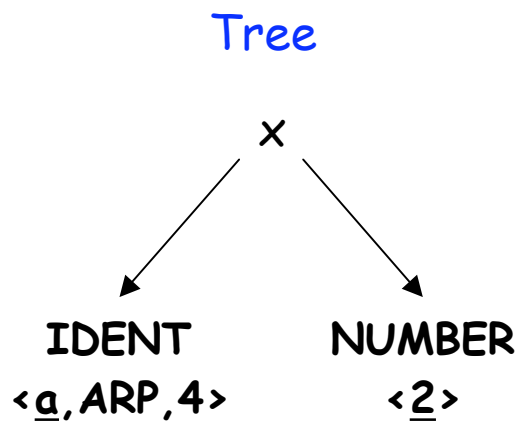
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Our treewalk code generator (Lec. 24) ran quickly

How good was the code?



Treewalk Code

```
loadI 4 ⇒ r5
loadAO r0, r5 ⇒ r6
loadI 2 ⇒ r7
mult r6, r7 ⇒ r8
```

Desired Code

```
loadAI r0, 4 ⇒ r5
multI r5, 2 ⇒ r7
```



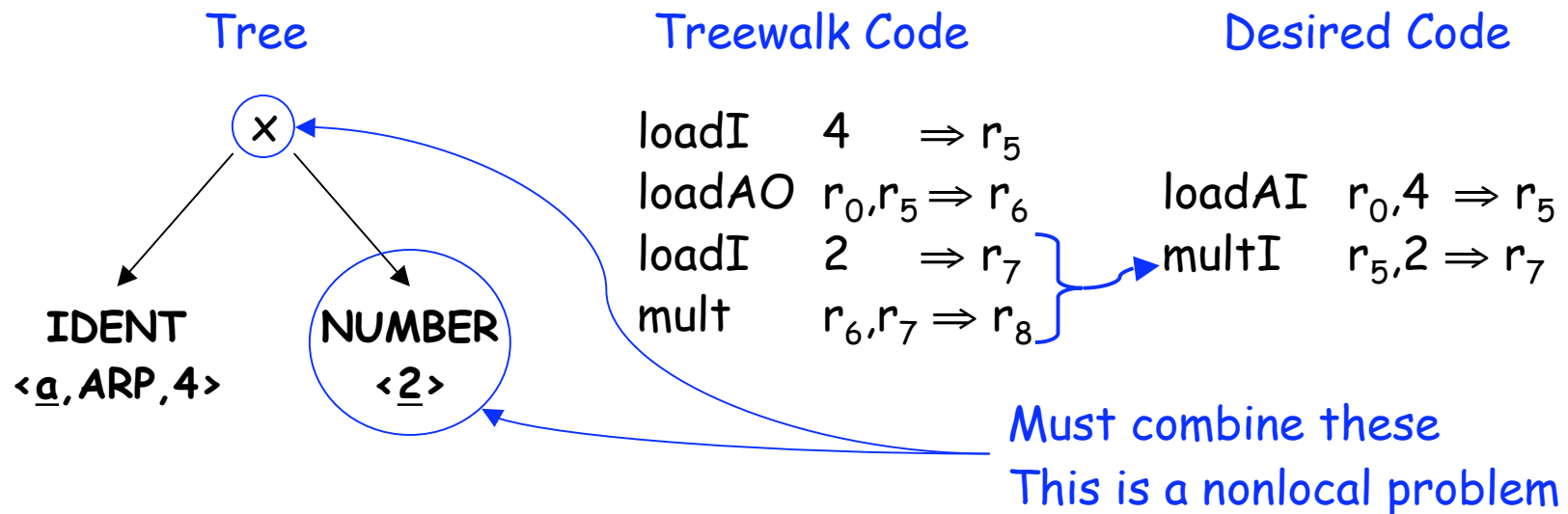

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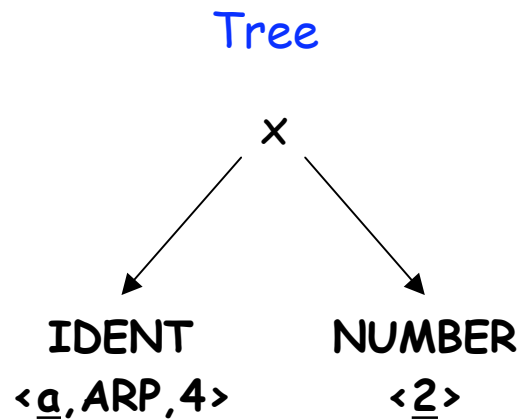
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Treewalk Code

```

loadI  4  => r5
loadAO r0,r5 => r6
loadI  2  => r7
mult   r6,r7 => r8
  
```

Desired Code

```

loadAI r0,4 => r5
add    r5,r5 => r7
  
```

(Note: The 'add' instruction and its register mapping are circled in blue in the original image, with an arrow pointing to the explanatory text below.)

Another possibility that might take less time & power (algebraic equivalence)



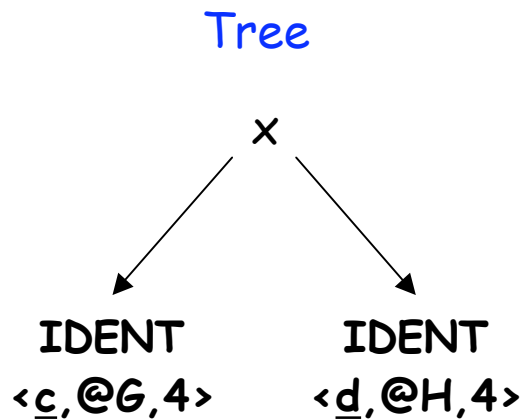
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator (Lec. 24) ran quickly

How good was the code?



Treewalk Code

```

loadI  @G => r5
loadI  4  => r6
loadAO r5,r6 => r7
loadI  @H => r7
loadI  4  => r8
loadAO r8,r9 => r10
mult   r7,r10 => r11
  
```

Desired Code

```

loadI  4      => r5
loadAI r5,@G => r6
loadAI r5,@H => r7
mult   r6,r7 => r8
  
```



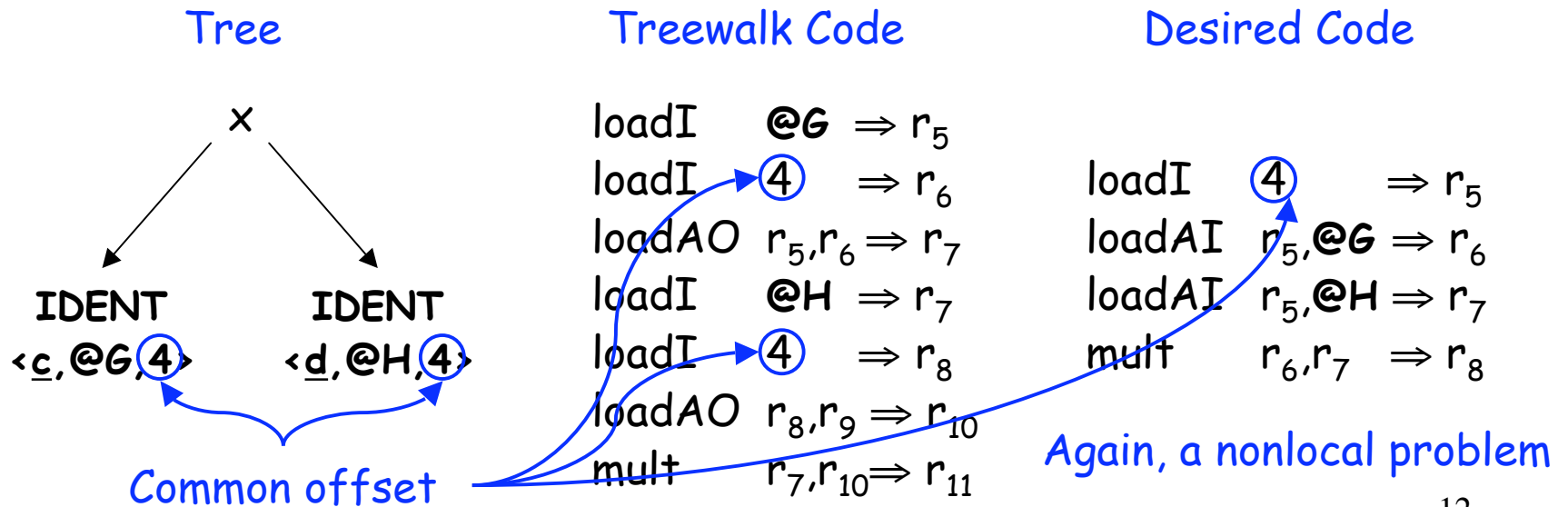
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator met the second criteria *(lec. 24)*

How did it do on the first ?



How do we perform this kind of matching ?



Tree-oriented IR suggests pattern matching on trees

- Process takes tree-patterns as input, matcher as output
- Each pattern maps to a target-machine instruction sequence
- Use dynamic programming or bottom-up rewrite systems

Linear IR suggests using some sort of string matching

- Process takes strings as input, matcher as output
- Each string maps to a target-machine instruction sequence
- Use text matching (Aho-Corasick) or peephole matching

In practice, both work well; matchers are quite different



Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load

Original code

```
storeAI r1    ⇒ r0,8  
loadAI  r0,8 ⇒ r15
```

Improved code

```
storeAI r1    ⇒ r0,8  
i2i     r1    ⇒ r15
```



Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load
- Simple algebraic identities

Original code

addI $r_2, 0 \Rightarrow r_7$
mult $r_4, r_7 \Rightarrow r_{10}$

multI $r_5, 2 \Rightarrow r_7$

Improved code

mult $r_4, r_2 \Rightarrow r_{10}$

add $r_2, r_2 \Rightarrow r_7$



Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load
- Simple algebraic identities
- Jump to a jump

Original code

```
    jumpI    → L10  
L10: jumpI    → L11
```

Improved code

```
L10: jumpI    → L11
```


Peephole Matching

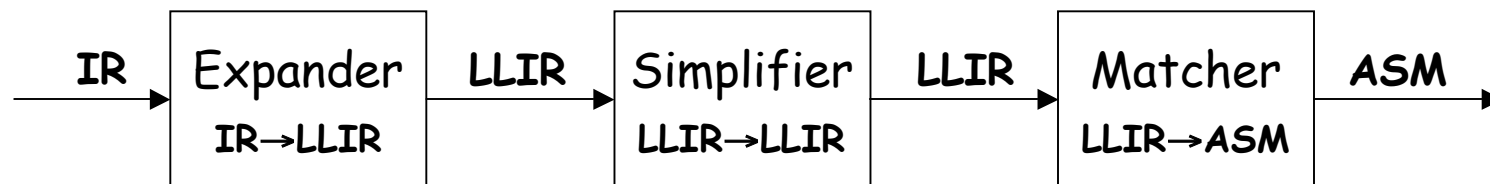


Implementing it

- Early systems used limited set of hand-coded patterns
- Window size ensured quick processing $O(n^2) \Rightarrow O(n)$

Modern peephole instruction selectors *(Davidson)*

- Break problem into three tasks



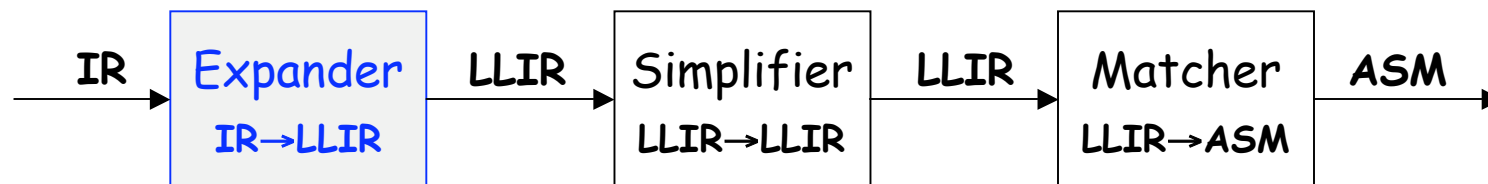
- Apply symbolic interpretation & simplification systematically

Peephole Matching



Expander

- Turns IR code into a low-level IR (LLIR) such as RTL
- Operation-by-operation, template-driven rewriting
- LLIR form includes all direct effects *(e.g., setting cc)*
- Significant, albeit constant, expansion of size



Peephole Matching



Simplifier

- Looks at LLIR through window and rewrites it
- Uses forward substitution, algebraic simplification, local constant propagation, and dead-effect elimination
- Performs local optimization within window



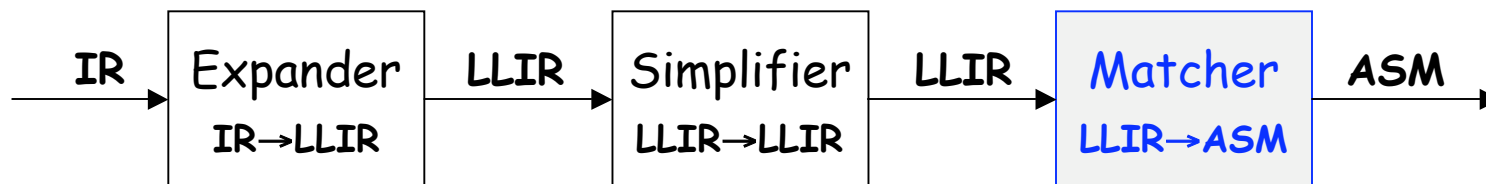
- This is the heart of the peephole system
 - Benefit of peephole optimization shows up in this step

Peephole Matching



Matcher

- Compares simplified LLIR against a library of patterns
- Picks low-cost pattern that captures effects
- Must preserve LLIR effects, may add new ones (*e.g., set cc*)
- Generates the assembly code output



Example



$x - 2 \times y$ *becomes*

Original IR Code

OP	Arg ₁	Arg ₂	Result
mult	2	y	t ₁
sub	x	t ₁	w

Symbolic names for memory-bound variables

Example



$x - 2 \times y$ becomes

Original IR Code

OP	Arg ₁	Arg ₂	Result
mult	2	y	t ₁
sub	x	t ₁	w

Expand



Symbolic names for memory-bound variables

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Example



LLIR Code

```
r10 ← 2
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r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(r0 + @y)
r14 ← 2 × r13
r17 ← MEM(r0 + @x)
r18 ← r17 - r14
MEM(r0 + @w) ← r18
```

Example



LLIR Code

$r_{13} \leftarrow \text{MEM}(r_0 + @y)$

$r_{14} \leftarrow 2 \times r_{13}$

$r_{17} \leftarrow \text{MEM}(r_0 + @x)$

$r_{18} \leftarrow r_{17} - r_{14}$

$\text{MEM}(r_0 + @w) \leftarrow r_{18}$

Match



Iloc Code

loadAI $r_0, @y \Rightarrow r_{13}$

multI $2 \times r_{13} \Rightarrow r_{14}$

loadAI $r_0, @x \Rightarrow r_{17}$

sub $r_{17} - r_{14} \Rightarrow r_{18}$

storeAI $r_{18} \Rightarrow r_0, @w$

- Introduced all memory operations & temporary names
- Turned out pretty good code

Steps of the Simplifier

(3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
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r20 ← r0 + r19
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```

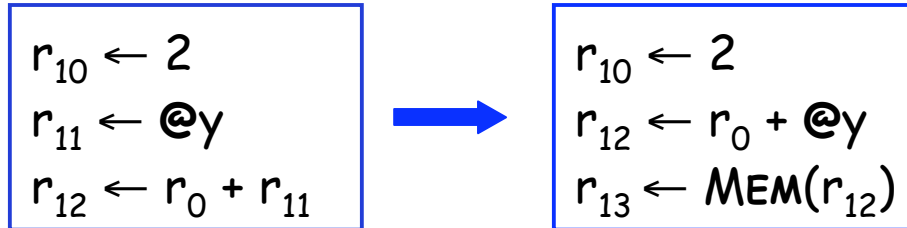
```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
```

Steps of the Simplifier (3-operation window)



LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow r_0 + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} \times r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow r_0 + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow r_0 + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

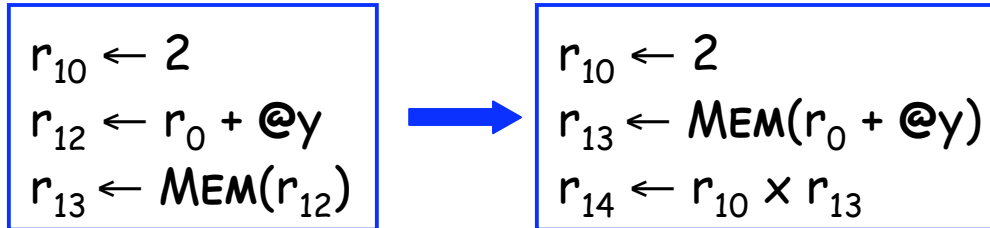


Steps of the Simplifier (3-operation window)



LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow r_0 + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} \times r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow r_0 + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow r_0 + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$



Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

```
r10 ← 2  
r13 ← MEM(r0 + @y)  
r14 ← r10 × r13
```



```
r13 ← MEM(r0 + @y)  
r14 ← 2 × r13  
r15 ← @x
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

```
r13 ← MEM(r0 + @y)
r14 ← 2 × r13
r15 ← @x
```

1st op it has rolled out of window



```
r14 ← 2 × r13
r15 ← @x
r16 ← r0 + r15
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

```
r14 ← 2 × r13  
r15 ← @x  
r16 ← r0 + r15
```



```
r14 ← 2 × r13  
r16 ← r0 + @x  
r17 ← MEM(r16)
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
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r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

r₁₄ ← 2 × r₁₃
r₁₆ ← r₀ + @x
r₁₇ ← MEM(r₁₆)



r₁₄ ← 2 × r₁₃
r₁₇ ← MEM(r₀ + @x)
r₁₈ ← r₁₇ - r₁₄

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

$r_{14} \leftarrow 2 \times r_{13}$
 $r_{17} \leftarrow \text{MEM}(r_0 + @x)$
 $r_{18} \leftarrow r_{17} - r_{14}$

$r_{17} \leftarrow \text{MEM}(r_0 + @x)$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
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r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

```
r17 ← MEM(r0+@x)
r18 ← r17 - r14
r19 ← @w
```

```
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
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r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

```
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
```



```
r18 ← r17 - r14
r20 ← r0 + @w
MEM(r20) ← r18
```

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

$r_{18} \leftarrow r_{17} - r_{14}$
 $r_{20} \leftarrow r_0 + @w$
 $MEM(r_{20}) \leftarrow r_{18}$



$r_{18} \leftarrow r_{17} - r_{14}$
 $MEM(r_0 + @w) \leftarrow r_{18}$

Steps of the Simplifier (3-operation window)



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

```
r18 ← r17 - r14
r20 ← r0 + @w
MEM(r20) ← r18
```



```
r18 ← r17 - r14
MEM(r0 + @w) ← r18
```

Example



LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(r0 + @y)
r14 ← 2 × r13
r17 ← MEM(r0 + @x)
r18 ← r17 - r14
MEM(r0 + @w) ← r18
```

Making It All Work



Details

- LLIR is largely machine independent
- Target machine described as LLIR → ASM pattern
- Actual pattern matching
 - Use a hand-coded pattern matcher
 - Turn patterns into grammar & use LR parser
- Several important compilers use this technology
- It seems to produce good portable instruction selectors

(RTL)

(gcc)

(VPO)

Key strength appears to be late low-level optimization